Truly Non-blocking Writes

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HotStorage Workshop, 2011

Introduction

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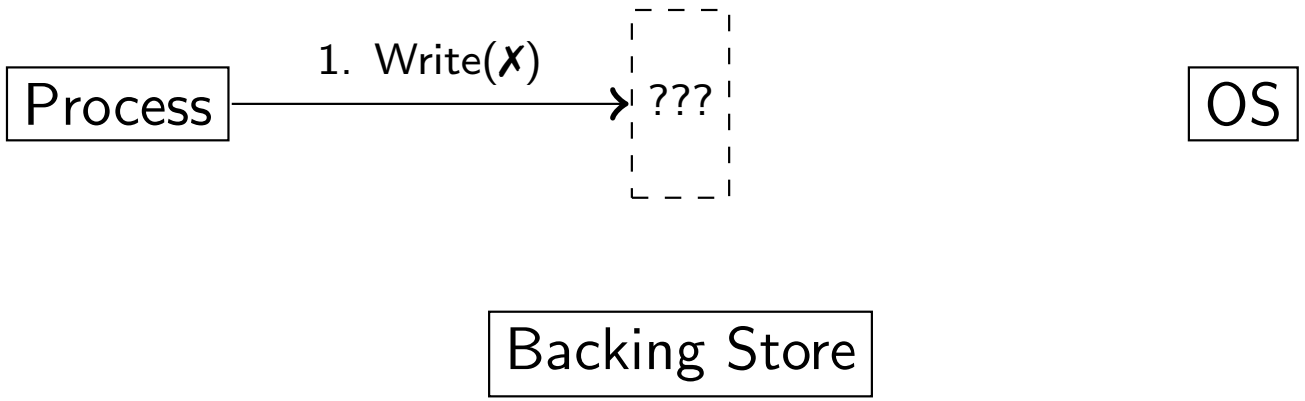

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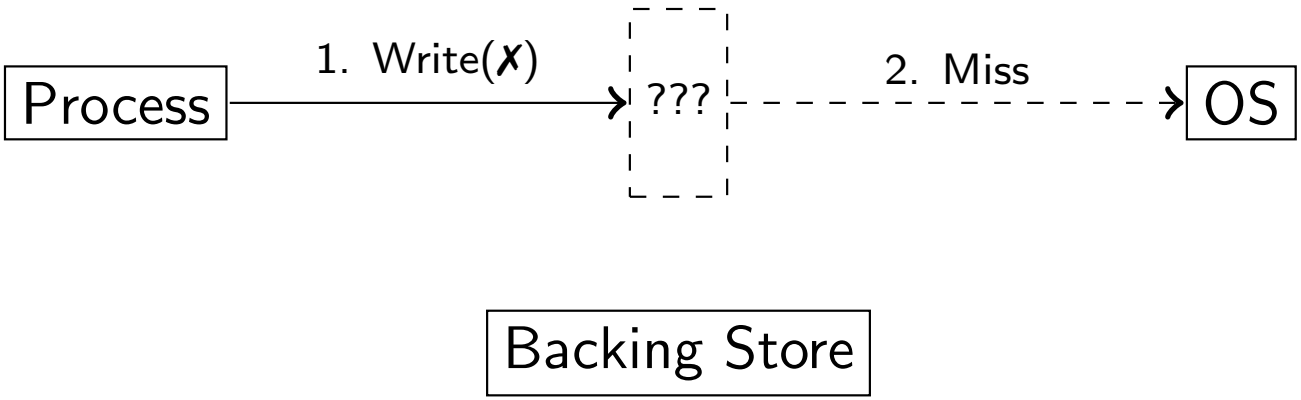

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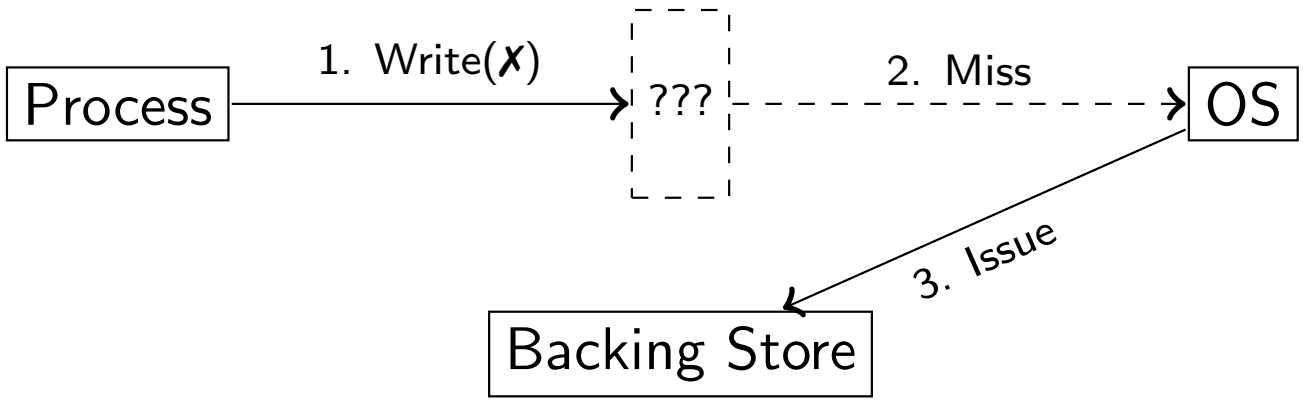

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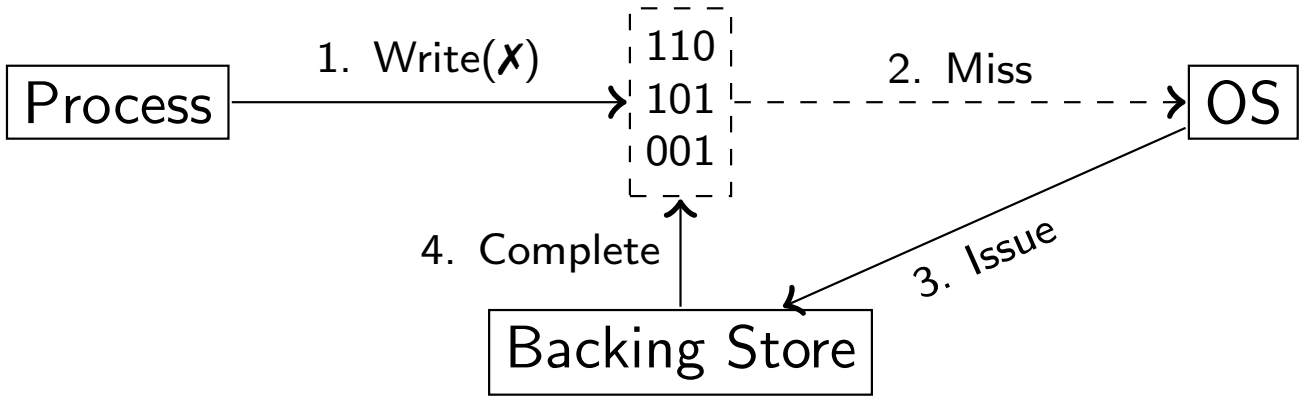

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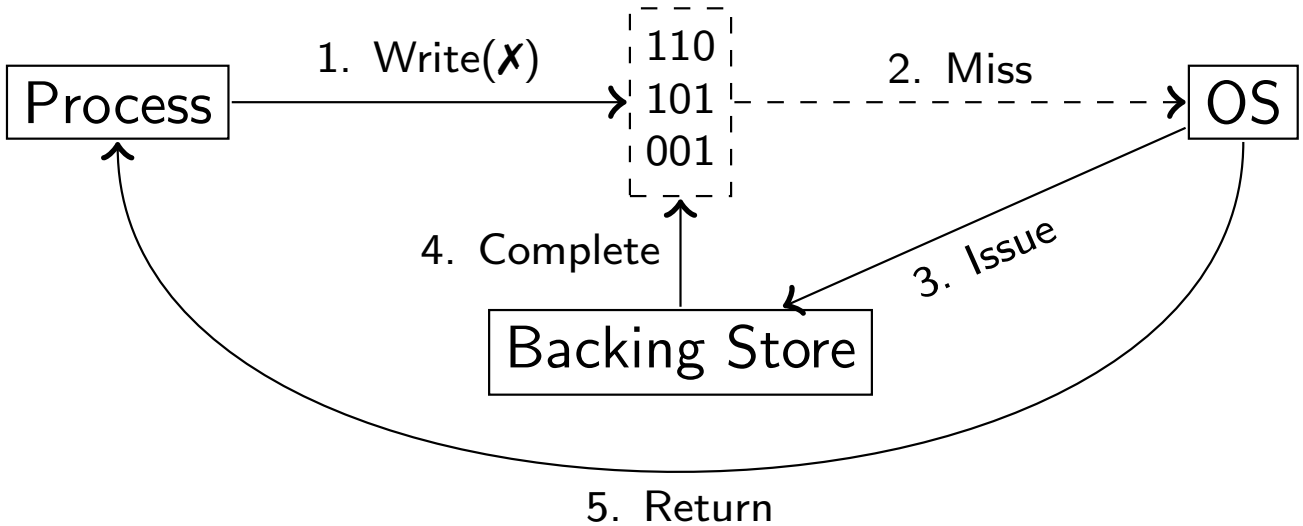

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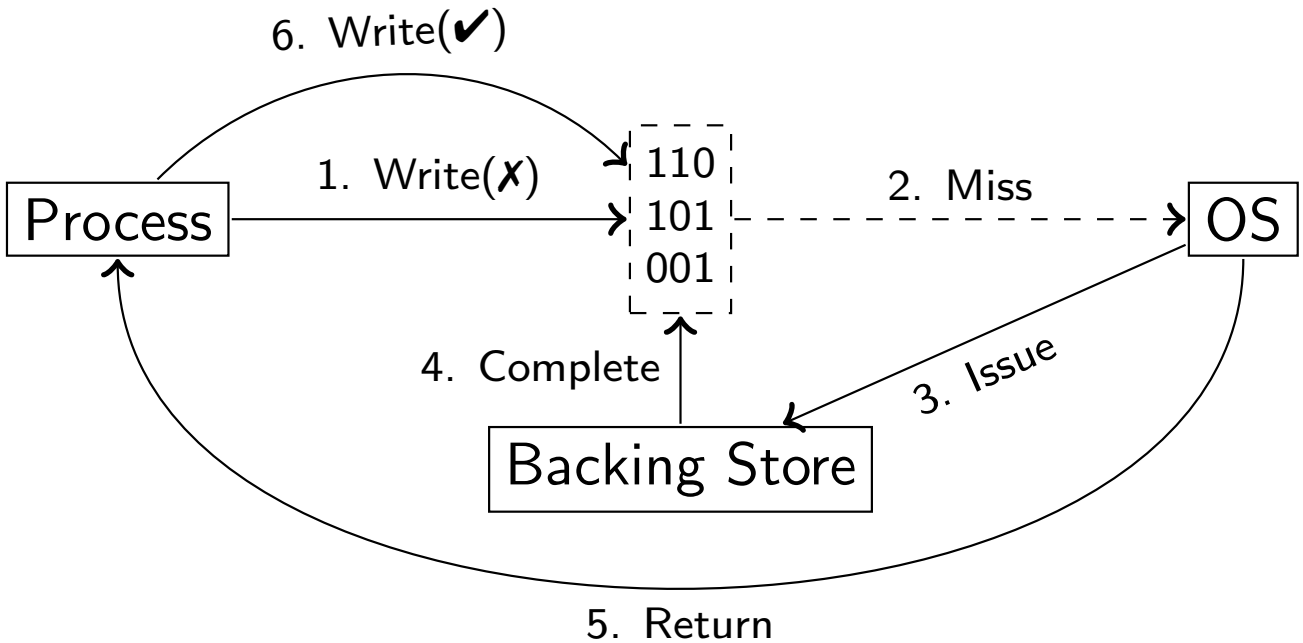

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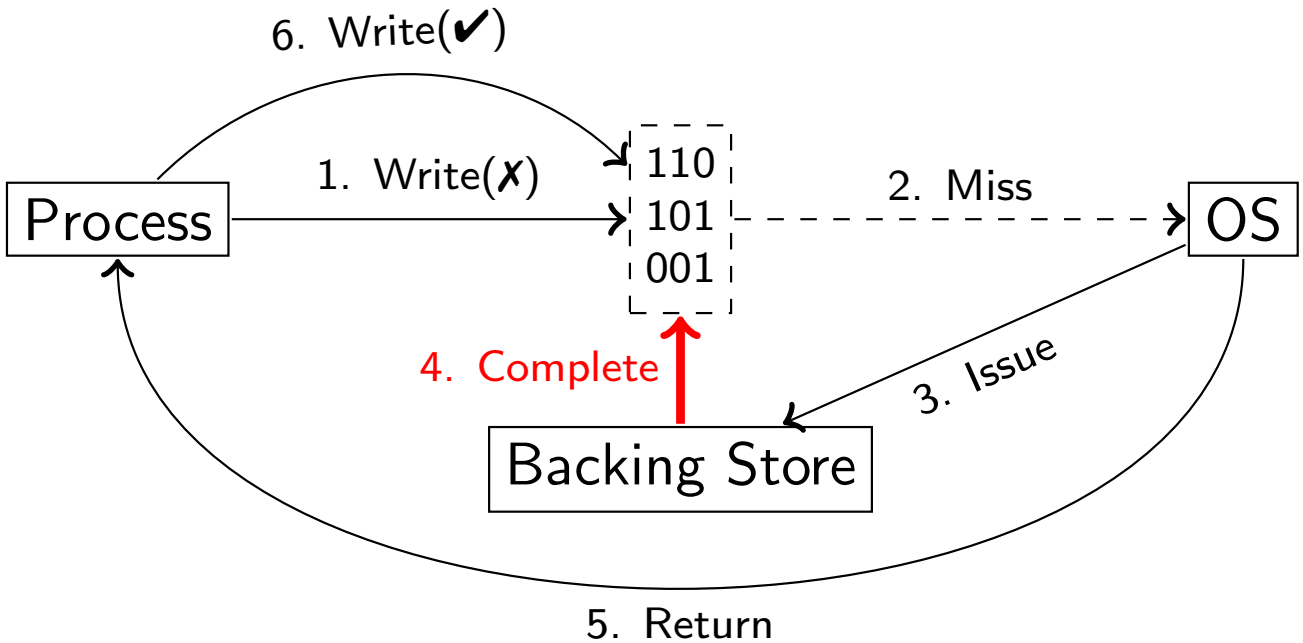

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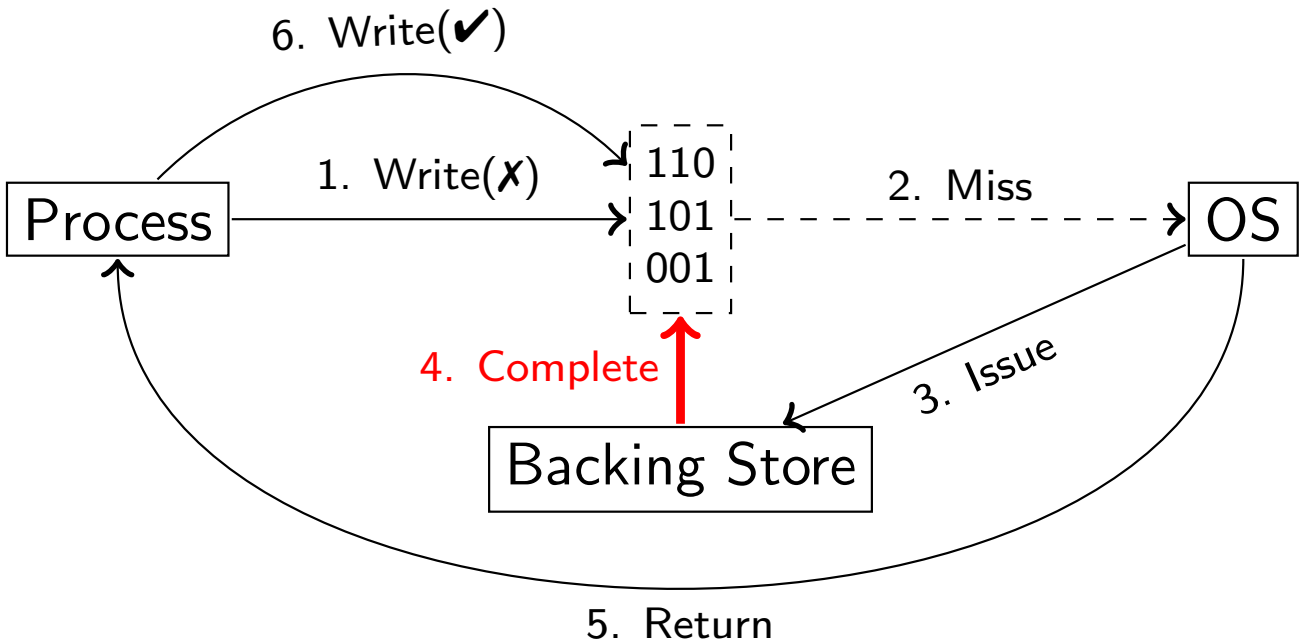

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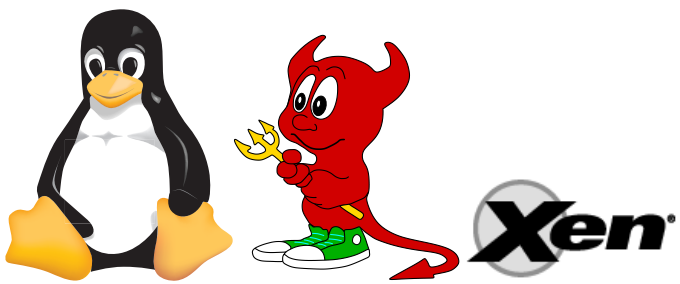

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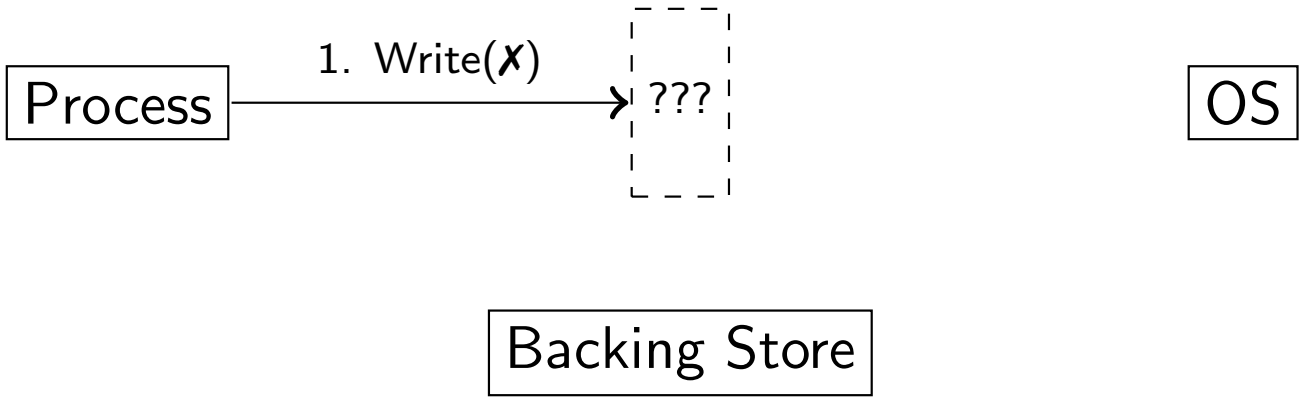

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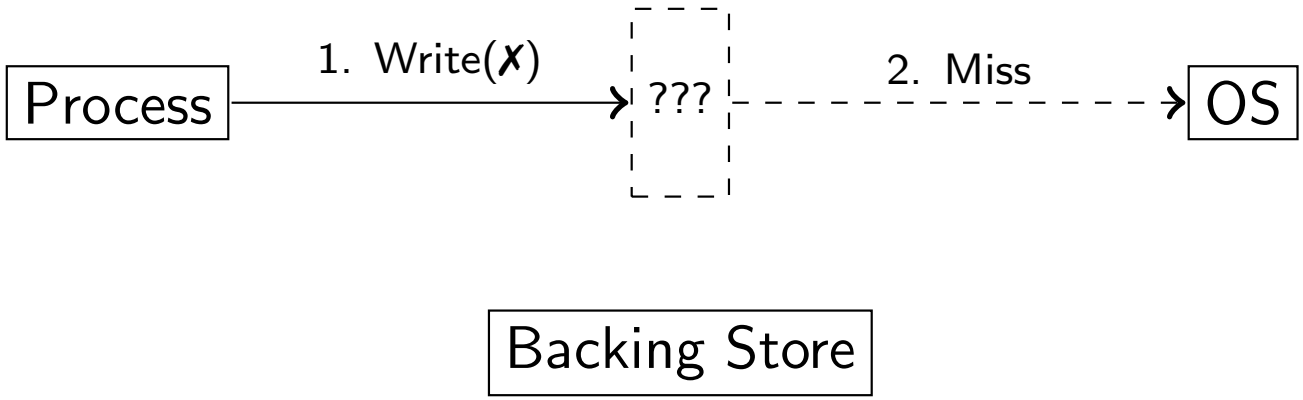

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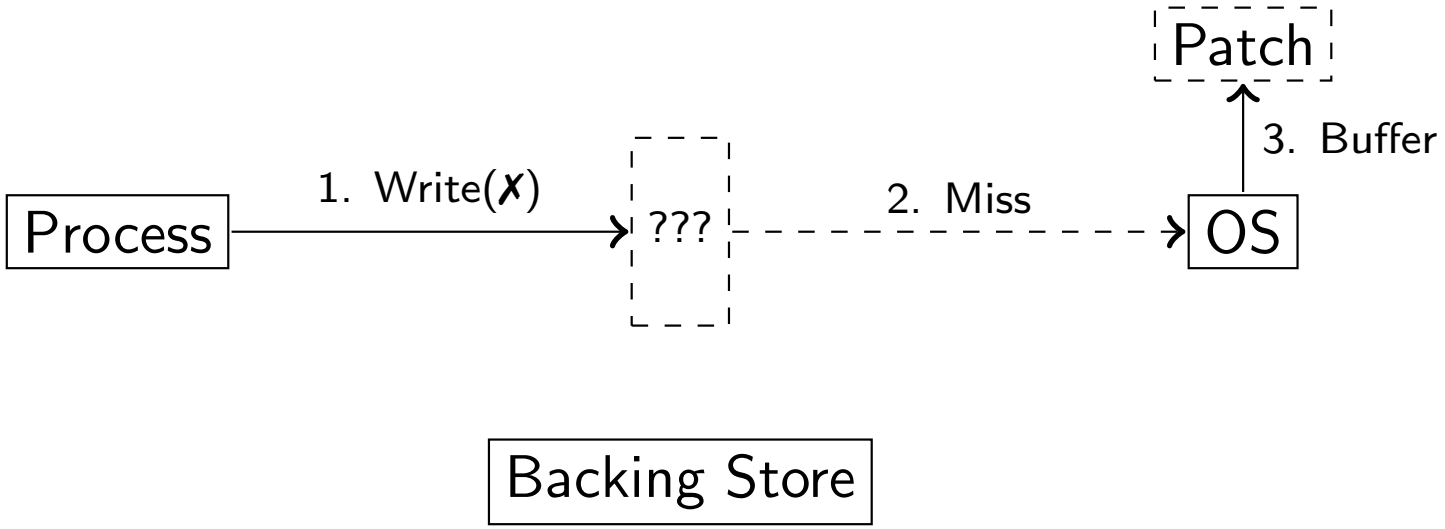

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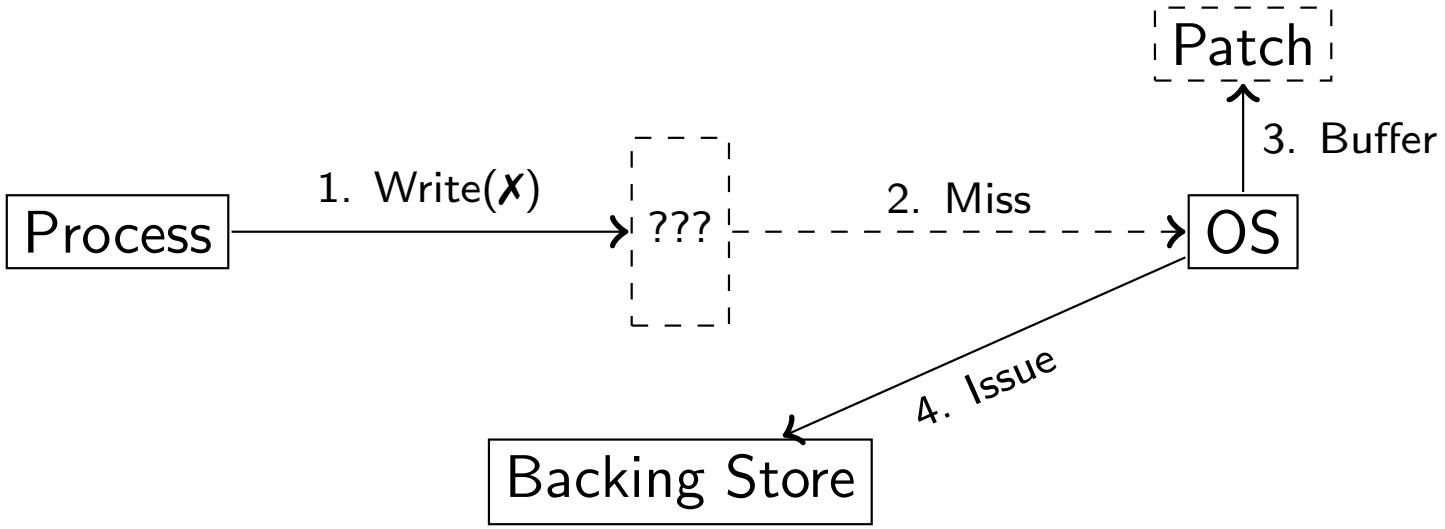

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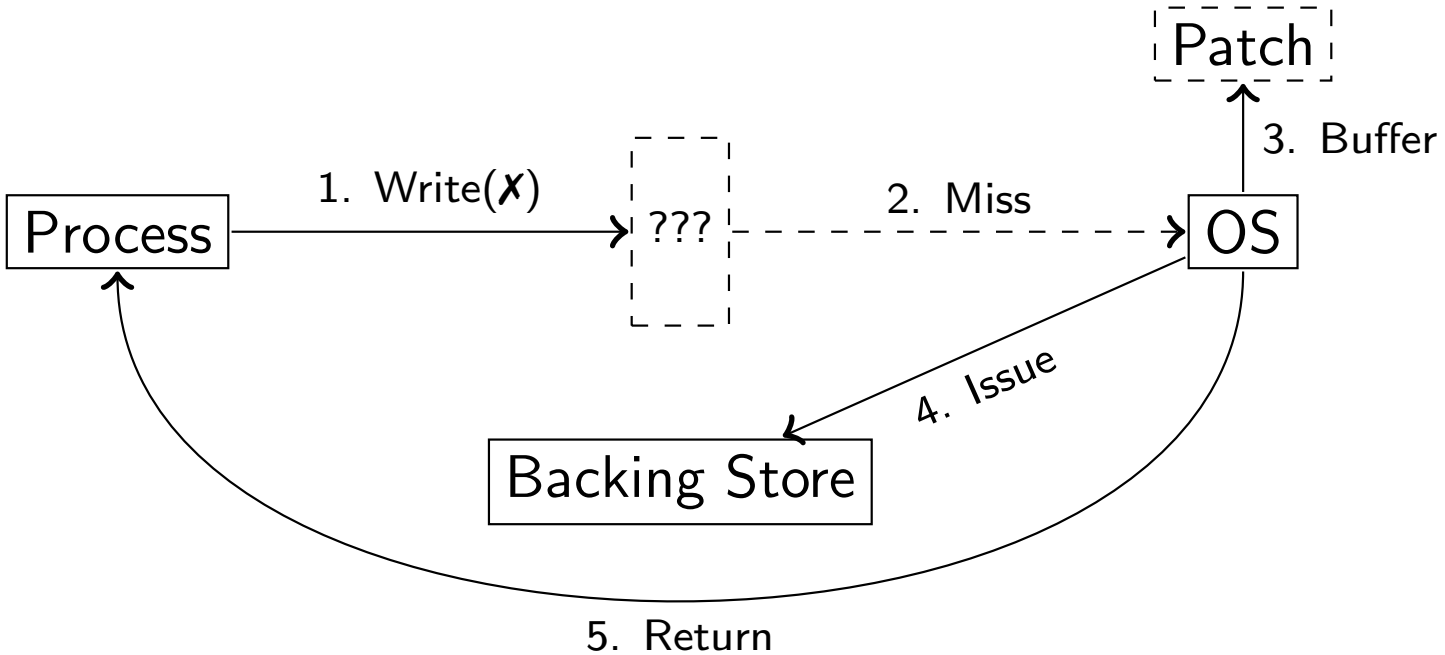

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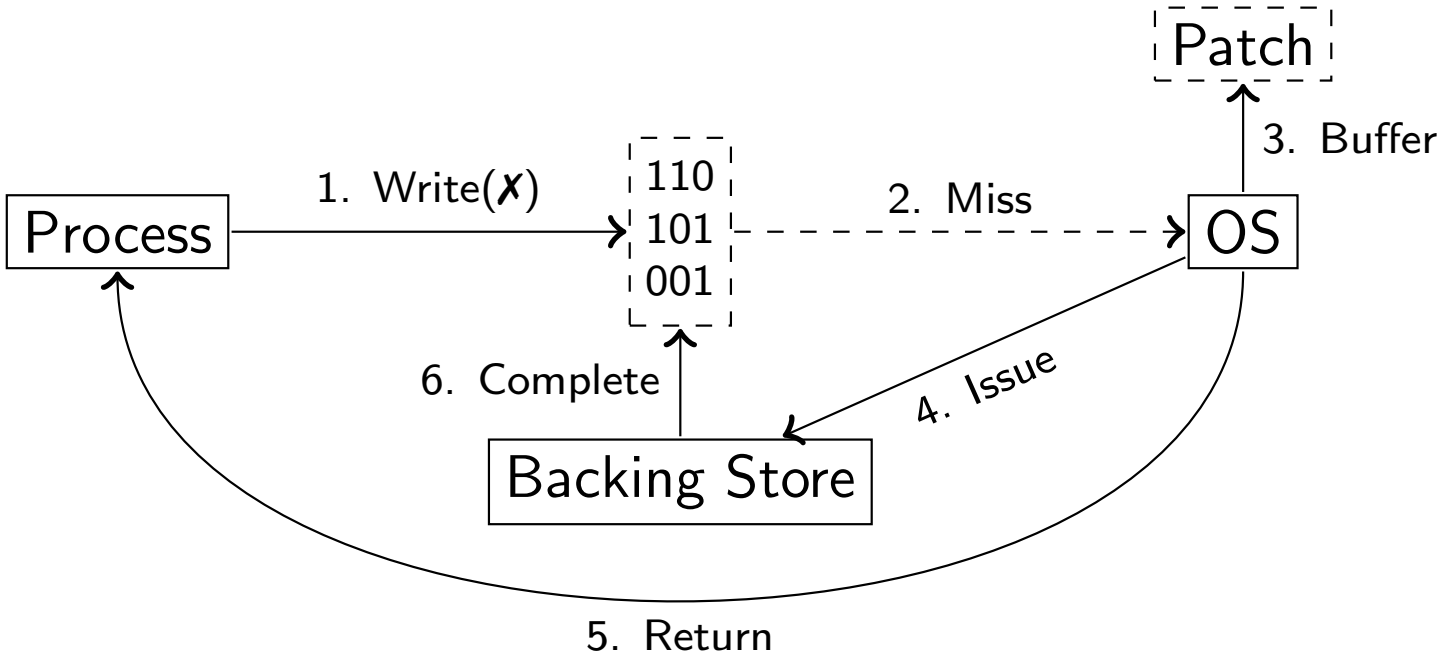

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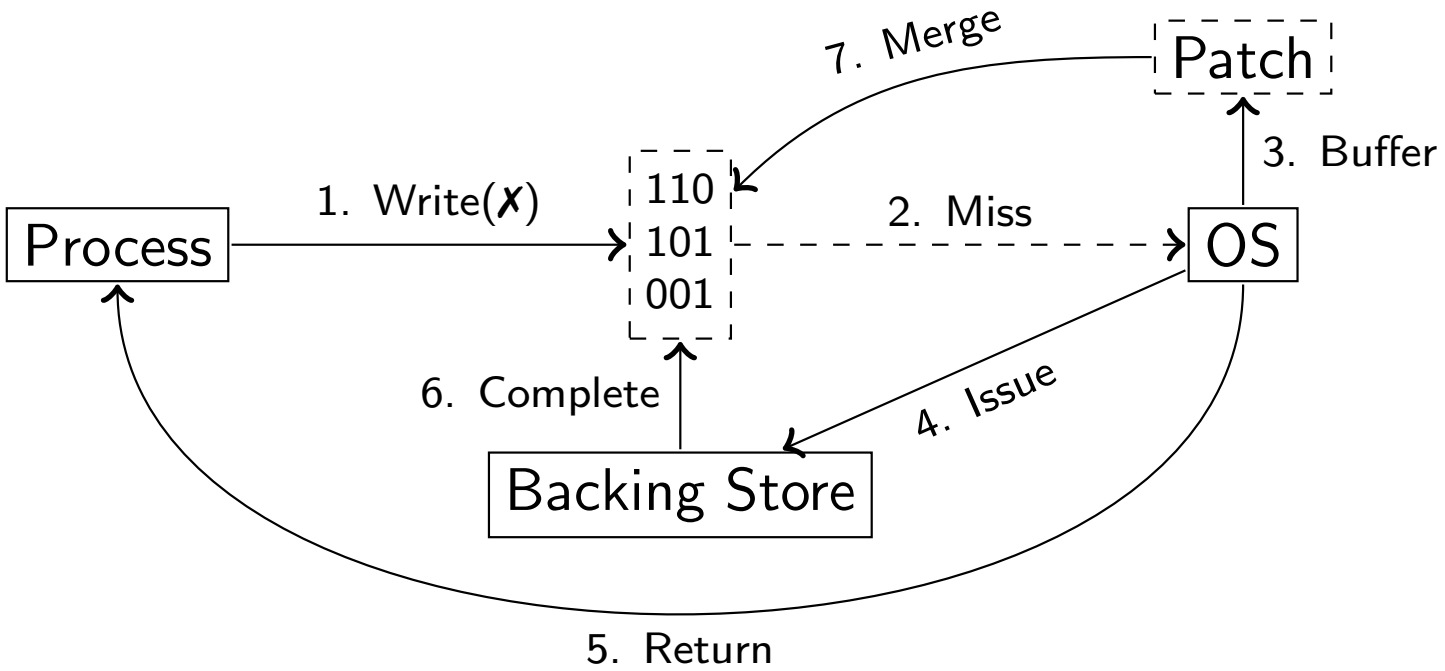

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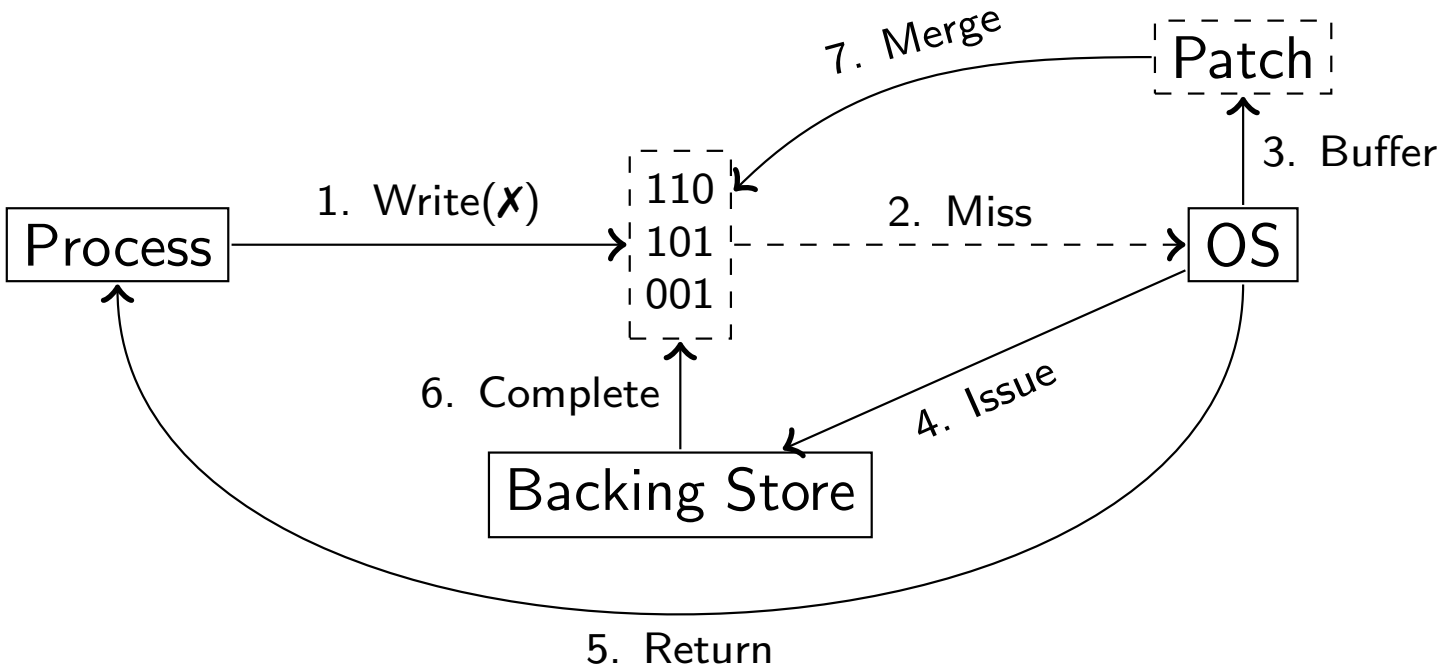
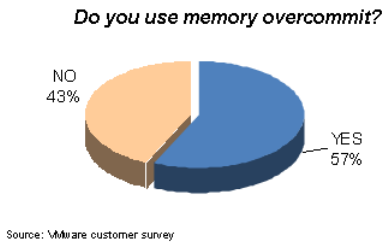


Benefits

1. Application execution time reduction
2. Increased backing store bandwidth usage

Motivation → Higher Fault Rates



Memory over-committed in virtualized environments

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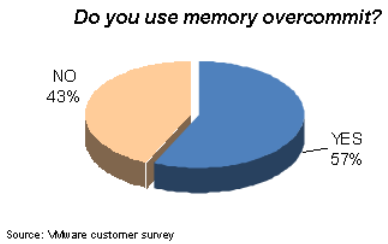

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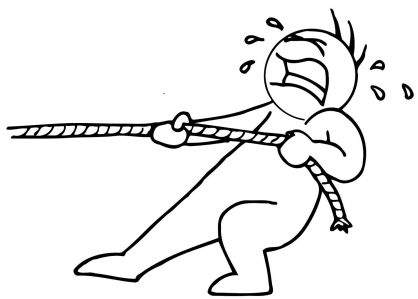

More process running with multi-core and virtualized environments

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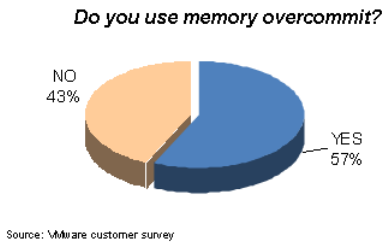

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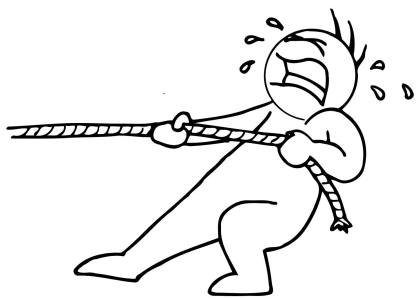

More process running with multi-core and virtualized environments

Memory hierarchy moving towards a more active and faster backing store

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- RAM set to 50% of app footprint

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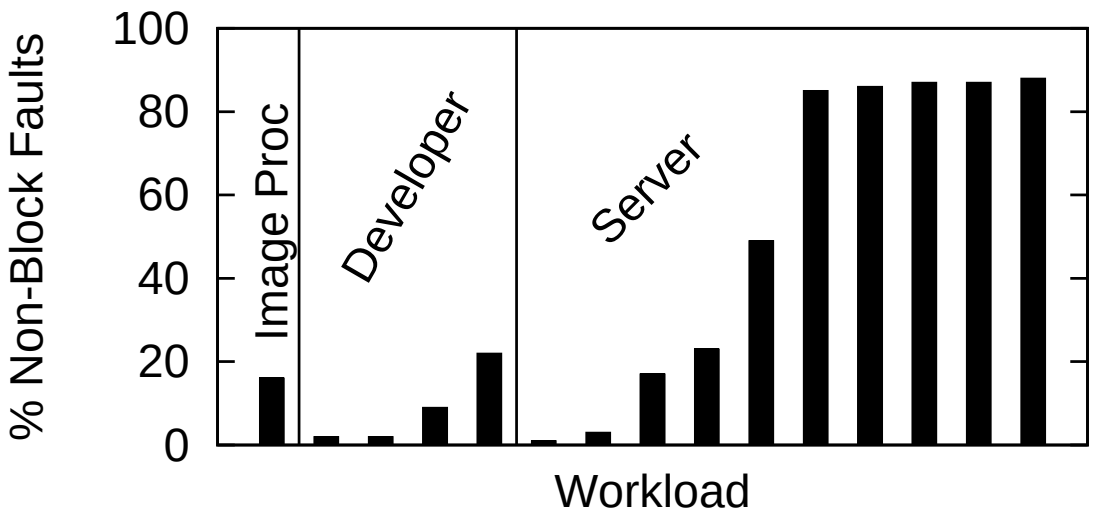
- RAM set to 50% of app footprint
- Up to 80% of page faults benefit



Related Work

Alternatives to non-blocking writes:

Perfect DRAM Provision

Unpredictable or unbounded.

Prefetching

Can incur false positives and false negatives.

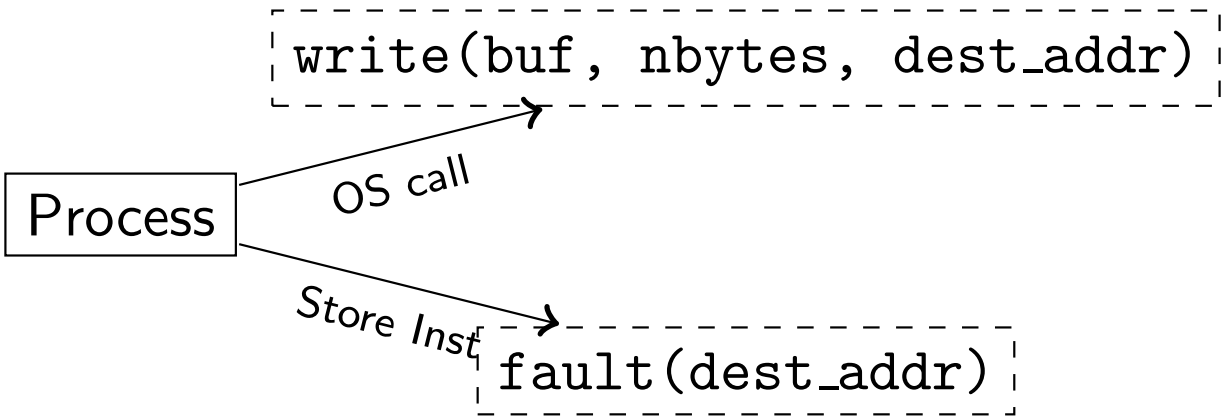
Asynchronous System Calls

1. Do not work with memory mapped pages
2. Written data *not* immediately available for reading

Solution Challenges

Process

Solution Challenges



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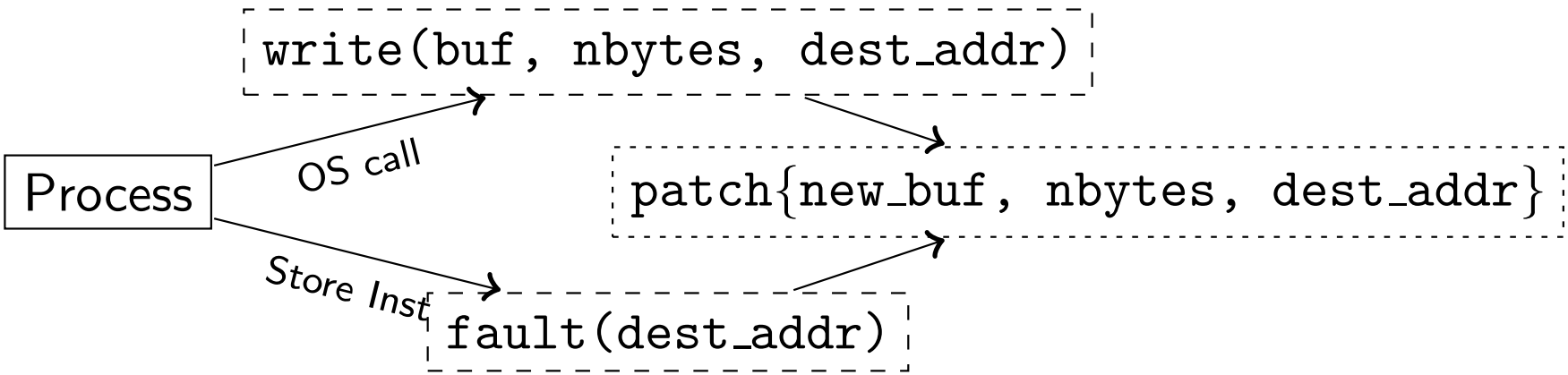

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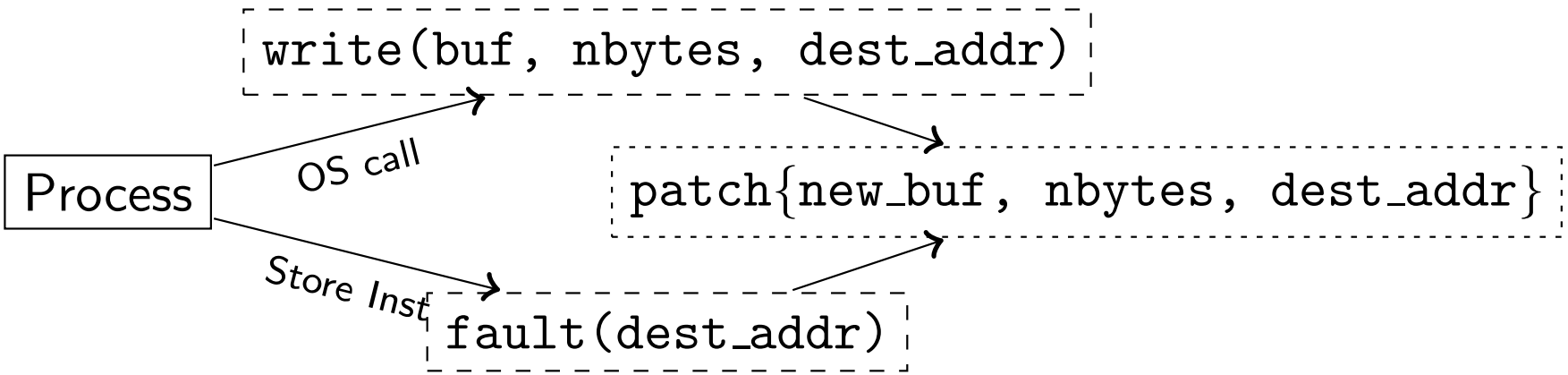

Information Per Non-blocking Write

Information
Write Offset
Data Written
Size of Data

Solution Challenges

Information Per Non-blocking Write

Information	Supervised `write()`
Write Offset	✔
Data Written	✔
Size of Data	✔

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Information	Supervised `write()`	Unsupervised Fault
Write Offset	✔	✔
Data Written	✔	✗
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Write Offset	✔	✔
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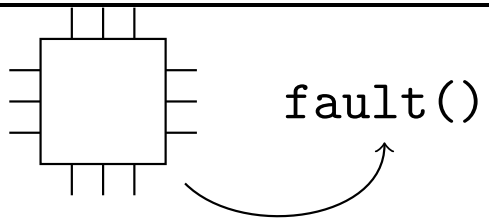
Handling Unsupervised Writes

Approach	Description	Fast	All Arch?	Low Mem?
Full Feature Hardware		✔	✗	✔

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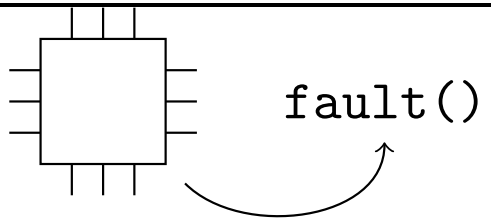
Full Feature Hardware		✔	✗	✔
Opcode Disassembly	4 bytes offset `sw $t1, 0xff` data	✔	✗	✔

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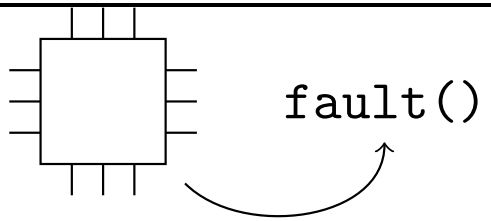
Full Feature Hardware		✔	✗	✔
Opcode Disassembly	4 bytes offset `sw $t1, 0xff` data	✔	✗	✔
Page Diff-Merge	Disk Page and 0-buffer or 1-buffer Updated Page	✗	✔	✗

Quantifying Benefits

1. Fraction of non-blocking write faults ✔
2. Outstanding write faults (over time)
3. Savings in execution time (new!)

Virtual Memory Simulator

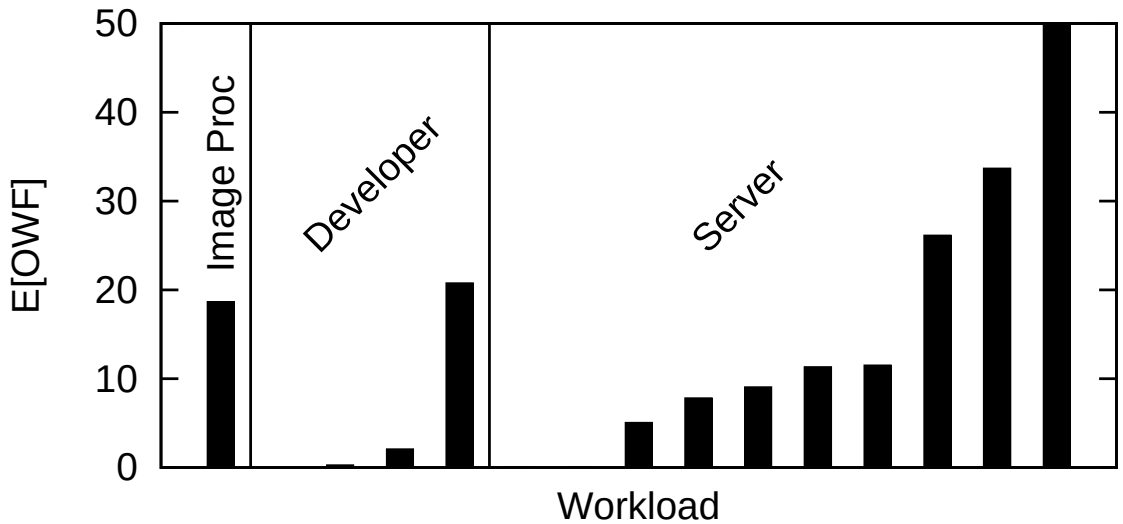
Input RAM size & Full System Memory Traces

Output Performance statistics

- Memory size set to 50% of workloads footprint
- Creating patches is not required

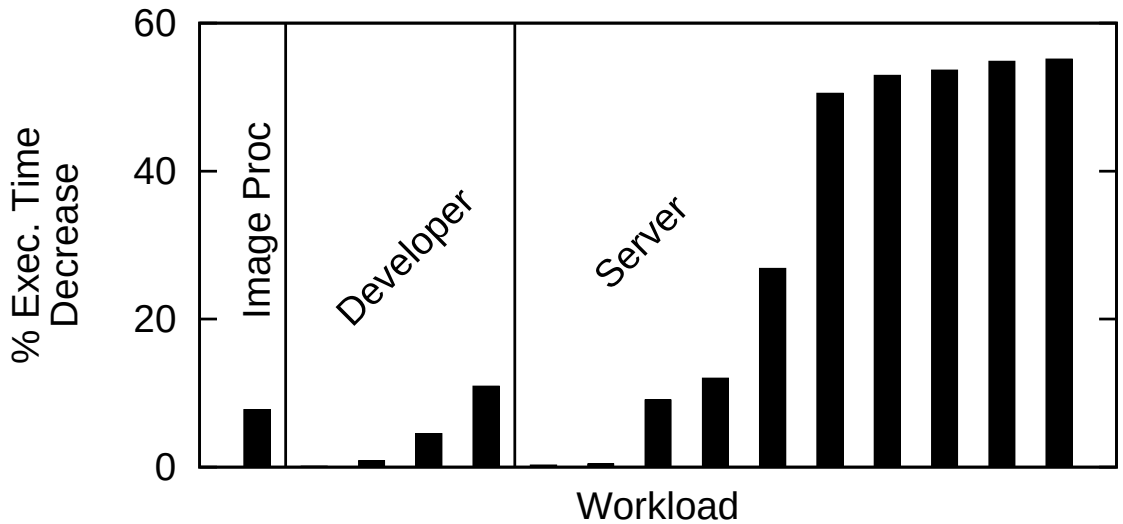
Quantifying Benefits → Metric

- How to measure the additional parallelism?
- **Outstanding Write Faults (OWF):** # of parallel write faults at any time
  - ✓ OWF $\leq$ OIO
  - ✓ OWF ≤ 1 for single threaded applications
  - ✓ OWF ≥ 0 when using non-blocking writes
- We need the variations over time as well
- **E[**OWF**]**: time-weighted average OWF



Quantifying Benefits → Time Reduction

- These results are not in the paper
- Execution time = Trace time + Synchronous read time
- Write time of dirty page on evictions ignored
- Rough estimate: error proportional to the number of dirty pages evicted



Conclusions and Future Work

- We presented non-blocking writes: a technique to eliminate read-before-writes
  - ✓ Reduced execution time
  - ✓ Increased device usage
- We estimate a reduction times of 0.1-54%
- In the future, we are planning to implement non-blocking writes to better study its implications
  - ✓ What workloads benefit from Non-blocking writes?

Questions?

Virtual Memory Simulator

Input: RAM size & Mem Traces

Output: Per Entry: Timestamp and event (hit, miss, evict);
Global: Performance stats.

- Writes to out-of-core pages considered non-blocking
- Non-blocking status revoked when:
  1. The page is read before I/O completion
  2. The page is evicted before I/O completion

Quantifying Benefits → Full System Memory Traces

Modified x86 software-MMU QEMU to log all memory accesses:

- *Instruction count, CR3, virtual/physical address, access-mode, page privileges.*

Workloads

Type	#	Footprint Avg/Std (MB)
Server	10	294/158
Developer	4	269/183
Image	1	149/0

Solution Approaches → Page Diff-Merge

1. Write in two pages: 0-page and 1-page.
2. Merge with *and* and *or*.

Process

Backing Store

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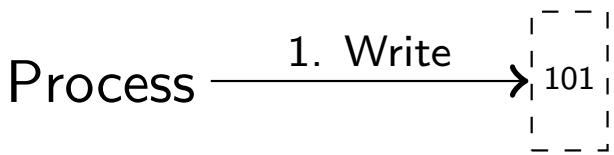

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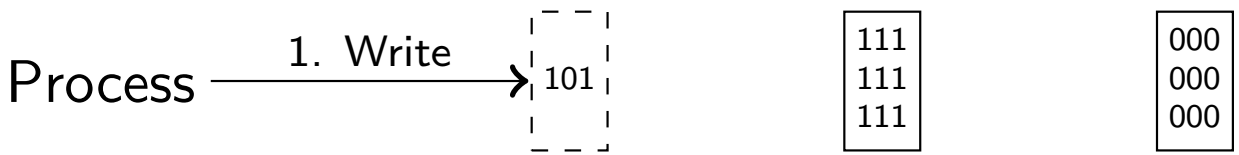

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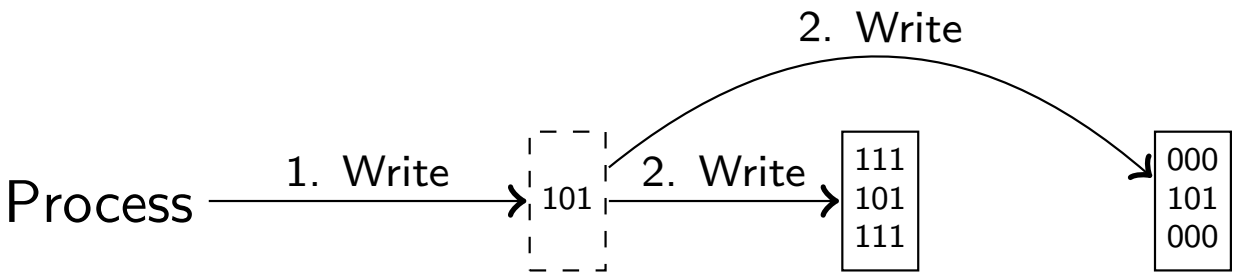

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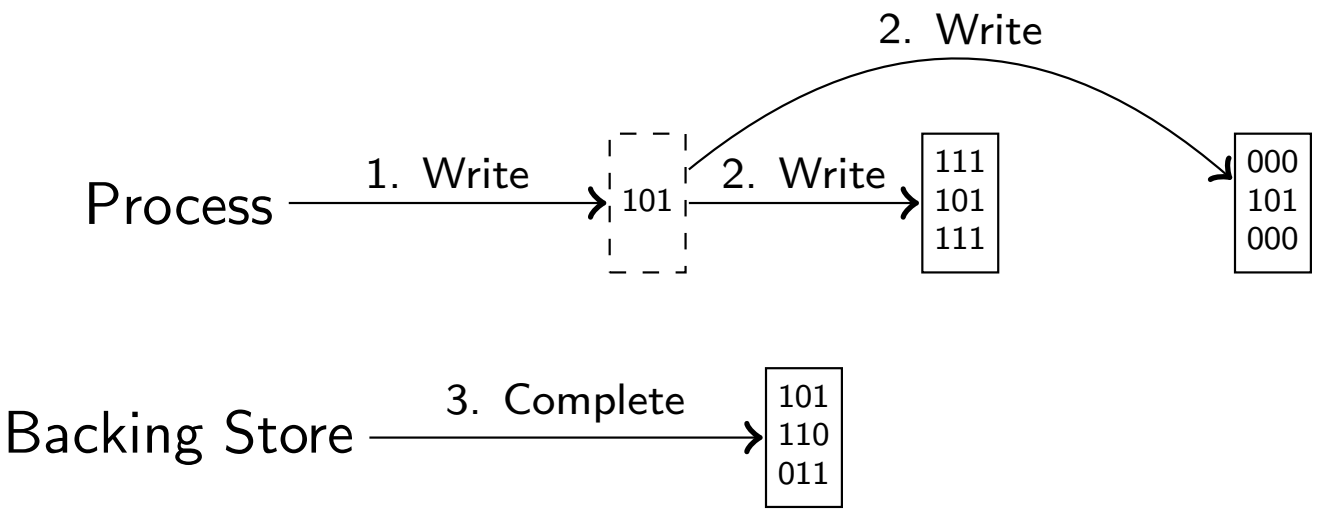

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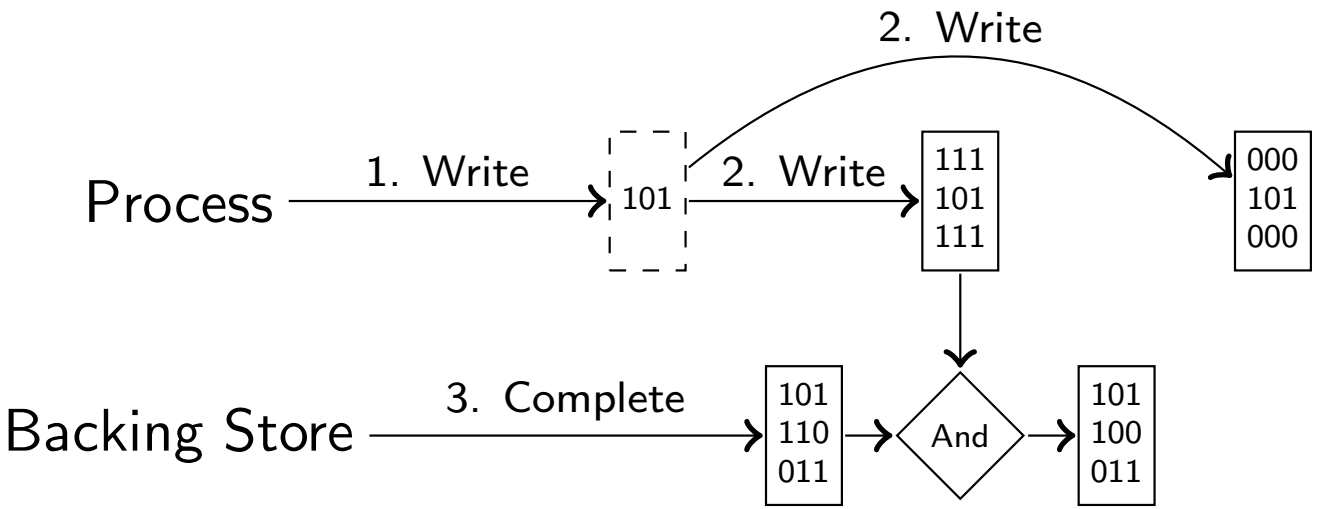

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